A New Maxmin Approach to Solve Time Minimization Transportation Problem with Mixed Constraints

Research Article

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#### Abstract

Some emergency cases transportation of goods is essential than transportation cost. Regarding the issue this paper presents an efficient algorithm to take full advantage of the transported unit with minimum time pass over the shipping cost in which the problem consists of equality and inequality constraints. Furthermore, numerical examples are used to demonstrate the method. A numerical example demonstrates how our suggested algorithm is very simple to comprehend and gets the most shipping products per unit.


Keywords: Transportation problem, Time minimization Transportation problem, Mixed constraints, Maximum flow, Minimum time

## 1. Introduction

Transportation problem (TP) is a particular type of linear programming issue, and closely related to moving goods from one place to another. The goal is to determine the minimal amount of good that has to be delivered from each source to a particular destination while also meeting all supply and demand constraints for the relevant sources and destinations. It is argued that there are equality restrictions on a balanced transportation problem if the overall supply and demand values are equal.
French mathematician Monge (1781) formalized the transportation issue. Russian mathematician and economist L. Kantorovich (1958) made significant advancements in the discipline during World War
II. Hitchcock (1941) was the one who originally presented the issue in its classic form. From the first introduction, different efficient methods of solution illustrated by Koopman (1947), Dantzig (1951) and then Charnes et al. (1953).

The total transportation problem is converted to Transportation Problem with Mixed Constraints (TPMC) if the supply of a source is significantly extended or lowered and the demand of a destination is similarly substantially increased or diminished. For More clarification, In the case of a general transportation issue, a factory could be required to produce 20 units of a product, but owing to an unforeseen circumstance, it might be able to produce fewer or more than 20 units instead. Similar to this, when a warehouse needed 40 pieces

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of inventory, they unexpectedly needed more or fewer units. That is why there is a transportation issue with several limits. Because of this, the transportation issue differs little from the balanced transportation issue. As a result, in real life, the majority of transportation issues include many limitations.
Brigden (1974) first carefully examined a unique sort of TPMC. Numerous research work on TPMC are found in the literature of Klingman et al. (1974), Isermann (1982), Adlakha et al. (2000, 2006), Pandian et al. (2010a-2010d), Mondal et al. (2012, 2014, 2015), Akilbasha et al.(2017).
The primary objective shouldn't be to cut transportation costs. In reality, a crisis situation might arise when the cost of expanding a moved unit is comparable. We would want to describe this urgent situation: itis very essential to send disaster relief to an area that is affacted by a sudden hurricane, earthquak, tsunami or flood. Some examples are mountain collaps, building collapse like Rana Plaza collaps or a daring rescue mission, suddenly happened earthquake in Turkey etc. When a medical emergency arises (such as corona, diarrhea, typhoid, dengue, chikungunya, or any other disease in epidemic form), enough medication or food should be sent to the affected region as quickly as feasible.
The best disaster assistance may readily compare to consider cost in this emergency situation. We provide a way to transfer a maximum number of units while meeting all supply and demand constraints in the shortest amount of time feasible without considering the cost of transportation.So, Time Minimization Transportation Problem arise.
Time Minimization Transportation Problem (TMTP) with equality constraints addressed by Peter L. Hammer for the first time in 1969. Then day by day many research worked found such as Garfinkel and Rao (1971), Wlodzimierz Szwarc (1971).A good numbers of paper has already been published on TMTP with equality requirements such as:Bhatia (1977), Sharma, J.K. and Swarup, K., (1977), Pandian and Natarajan (2011), Uddin,
M.S., (2012), Nikolić, I., (2007),Md. Main Uddin et al. (2013), Mollah Mesbahuddin Ahmed (2014, 2015), Khan, A. R. et. al. (2016), Khan, A.R. (2017) have shown TMTPs and its variants in their research. They presenteddifferent algorithms for solving TMTP with equality constraints.
From this survey of the literature, we can see that while there has been a sizable amount of study into finding the best solution to the time-minimizing transportation problem considering equality constraints, relatively few studies have focused on the issue of mixed constraints. The Time Minimization Transportation Problem with Mixed Constraints (TMTP-MC) is extensively useful in many real life situations where distinct sources and destinations require variableamounts of supply and demand commoditiesrespectively depending on to their availability and requirements.Agarwal, S., \& Sharma, S. (2014) as far I know (probably) was first introduced TMTP-MC. Agarwal, S., \& Sharma, S. $(2018,2020)$ finding IBFS of TMTPMC. Rashid, F. et al. (2021) proposed a method to solve TMTP-MC. Mentioned above literature review on TMTP-MC we see that all methods are focused on to find minimum time. But in this paper we focused our method to maximize transported unit along with minimum time.

## 2. Formulation of General Transportation Problem

We represent the arrangement of $m$ source points from which a good is transported by $O_{1}, O_{2}, \cdots \cdots, O_{m}$ and the arrangement of $n$ destination points by $D_{1}, D_{2}, \cdots \cdots, D_{n}$ to which the product is transported. Supply point $i$ can supply at most $S_{i}$ units $(i=1,2, \cdots \cdots, m)$ whereas demand point $j$ must receive at least $d_{j}$ units $(j=1,2, \cdots \cdots, n)$ of shipped good. $x_{i j}$ is the amount transported and $c_{i j}$ is the unit transportation cost between supply point $i$ to destination point $j$ .Bearing in mind the overhead symbolization, the general TP can be characterized by the organization as in the following Fig.1:


Fig. 1: Network Representation of Transportation Problem
Also depiction of TP by using matrix form as follows:
Table1: Matrix Representation of TP.

| Factory | Warehouse |  |  |  | Supply |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | ... | n |  |
| 1 | $\mathrm{X}_{11}$ | $\mathrm{X}_{12}$ | $\cdots$ | $\mathrm{X}_{1 \mathrm{n}}$ | $\mathrm{S}_{1}$ |
|  | $\mathrm{c}_{11}$ | $\mathrm{c}_{12}$ |  | $\mathrm{c}_{1 \mathrm{n}}$ |  |
| 2 | $\mathrm{X}_{21}$ |  | $\cdots$ | $\mathrm{X}_{2 \mathrm{n}}$ | $\mathrm{S}_{2}$ |
|  | $\mathrm{c}_{21}$ | $\mathrm{c}_{22}$ |  | $\mathrm{c}_{2 \mathrm{n}}$ |  |
| $\vdots$ | : | : | $\ldots$ | : | ; |
| m | $\mathrm{X}_{\mathrm{m} 1}$ | $\mathrm{x}_{\mathrm{m} 2}$ | $\ldots$ | $\mathrm{x}_{\mathrm{mn}}$ | $\mathrm{S}_{\mathrm{m}}$ |
|  | $\mathrm{c}_{\mathrm{m} 1}$ | $\mathrm{c}_{\mathrm{m} 2}$ |  | $\mathrm{c}_{\mathrm{mn}}$ |  |
| Demand | $\mathrm{d}_{1}$ | $\mathrm{d}_{2}$ | $\cdots$ | $\mathrm{d}_{\mathrm{n}}$ |  |

The TP is then expressed mathematically as For Balanced TP (1) becomes to: follows:

Minimize $\quad z=\sum_{i=1}^{m} \sum_{j=1}^{n} c_{i j} x_{i j}$.
Subject to $\sum_{j=1}^{n} x_{i j} \leq s_{i}$;
$i=1,2, \cdots \cdots, m$ (Supply constraints)
$\sum_{i=1}^{m} x_{i j} \geq d_{j} ; j=1,2, \cdots \cdots, n$ (Demand constraints)
$x_{i j} \geq 0 . \quad$ for all $i=1,2, \cdots \cdots, m$ and
$j=1,2, \cdots \cdots, n$.
$\operatorname{Minimize} \quad z=\sum_{i=1}^{m} \sum_{j=1}^{n} c_{i j} x_{i j}$.
Subject to $\sum_{j=1}^{n} x_{i j}=s_{i}$;
$i=1,2, \cdots \cdots, m$ (Supply constraints)
$\sum_{i=1}^{m} x_{i j}=d_{j} ; \quad j=1,2, \cdots \cdots, n$ (Demand
constraints)

$$
x_{i j} \geq 0 ; \quad \text { for all } i=1,2, \cdots \cdots, m
$$

and $j=1,2, \cdots \cdots, n$.
2.1. Formulation of TP with mixed constraints

Mathematically, a TPMC is defined as follows:
$\operatorname{Minimize} Z=\sum_{i=1}^{m} \sum_{j=1}^{n} c_{i j} x_{i j}$

Subject to Supply Constraints

$$
\begin{aligned}
& \sum_{j \in T} x_{i j}=a_{i}, \quad i \in S_{1} \\
& \sum_{j \in T} x_{i j} \geq a_{i}, \quad i \in S_{2} \\
& \sum_{j \in T} x_{i j} \leq a_{i}, \quad i \in S_{3}
\end{aligned}
$$

And $x_{i j} \geq 0$
Where $S_{1}, S_{2}, S_{3}$ are partitioned from $S, i \in S$, $i=1,2, \cdots \cdots, m$

And
Demand Constraints

$$
\begin{aligned}
& \sum_{i \in S} x_{i j}=b_{j}, \quad j \in T_{1} \\
& \sum_{i \in S} x_{i j} \geq b_{j}, \quad j \in T_{2} \\
& \sum_{i \in S} x_{i j} \leq b_{j}, \quad j \in T_{3}
\end{aligned}
$$

Where $T_{1}, T_{2}, T_{3}$ are partitioned from $T, j \in T$, $j=1,2, \cdots \cdots, n$

This is the standard form of TPMC.

### 2.2. Formulation of Time Minimization TP with Mixed constraints (TMTP-MC):

The structure of TMTP-MC is similar to the regular cost minimization problem with mixed constraints, except that the unit transportation cost $c_{i j}$ is replaced by the unit transportation time $t_{i j}$ required to effect a complete shipment of the goods from specific source $i$ to particular destination $j$.

Table : $\mathbf{2}$ Matrix Representation of TMTP-MC

| Factory | Warehouse |  |  |  | Supply |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | $\cdots$ | n |  |
| 1 | $\mathrm{X}_{11}$ | $\mathrm{X}_{12}$ |  | $\mathrm{X}_{1 \mathrm{n}}$ |  |
|  | $\mathrm{t}_{11}$ | $\mathrm{t}_{12}$ |  | $\mathrm{t}_{1 \mathrm{n}}$ |  |
| 2 | $\mathrm{X}_{21}$ | $\mathrm{X}_{22}$ | $\cdots$ | $\mathrm{X}_{2 \mathrm{n}}$ |  |
|  | $\mathrm{t}_{21}$ | $\mathrm{t}_{22}$ |  | $\mathrm{t}_{2 \mathrm{n}}$ |  |
| ! |  | ! | $\cdots$ | : | ; |
|  | $\mathrm{X}_{\mathrm{m} 1}$ | $\mathrm{x}_{\mathrm{m} 2}$ |  | $\mathrm{x}_{\mathrm{mn}}$ |  |
|  | $\mathrm{t}_{\mathrm{m} 1}$ | $\mathrm{t}_{\mathrm{m} 2}$ |  | $\mathrm{t}_{\mathrm{mn}}$ |  |
| Demand | $\geq 1=1 \leq \mathrm{b}_{1}$ | $\geq 1=1 \leq b_{2}$ | $\cdots$ | $\geq 1=1 \leq \mathrm{b}_{\mathrm{n}}$ |  |

$x_{i j}$ is the quantity shipped from source $i$ to demand point $j$ and $t_{i j}$ per unit transportation time from sources $i$ to demand point $j$.
$i=1,2, \cdots \cdots, m$ and $j=1,2, \cdots \cdots, n$
Therefore the LPP model of TMTP-MC is
Minimize Z: Max $t_{i j} \mid x_{i j}>0$
Subject to Supply Constraints

$$
\begin{aligned}
& \sum_{j \in T} x_{i j}=a_{i}, \quad i \in S_{1} \\
& \sum_{j \in T} x_{i j} \geq a_{i}, \quad i \in S_{2} \\
& \sum_{j \in T} x_{i j} \leq a_{i}, \quad i \in S_{3}
\end{aligned}
$$

And $x_{i j} \geq 0$
Where $S_{1}, S_{2}, S_{3}$ are partitioned from $S, i \in S$, $i=1,2, \cdots \cdots, m$

And
Demand Constraints

$$
\begin{aligned}
& \sum_{i \in S} x_{i j}=b_{j}, \quad j \in T_{1} \\
& \sum_{i \in S} x_{i j} \geq b_{j}, \quad j \in T_{2} \\
& \sum_{i \in S} x_{i j} \leq b_{j}, \quad j \in T_{3}
\end{aligned}
$$

Where $T_{1}, T_{2}, T_{3}$ are partitioned from $T, j \in T$, $j=1,2, \cdots \cdots, n$

This is the standard form of TMTP-MC.
Basic Assumptions:

1. $S_{2}$ must allocate at least $a_{i}$ units of supply which is equal to the total minimum demand and $T_{2}$ must receive at least $b_{j}$ units which is equal to the total minimum supply.
2. They start concurrently from their respective origins.
3. The time $t_{i j}$ is independent of the quantity $x_{i j}$.
4. Theoretical Development of Proposed Method.
Step-1: Formulation of TMTP-MC.
Step-2: For identifying total minimum supply and total minimum demand we set the condition below:
$a^{\prime}{ }_{i}=a_{i}, \boldsymbol{i} \in \boldsymbol{S}_{\mathbf{1}} \cup \boldsymbol{S}_{\mathbf{2}}$
$a_{i}^{\prime}=0, \quad i \in S_{3} \quad, i=1,2, \cdots \cdots, m$,
$j=1,2, \cdots \cdots, n$
and
$b^{\prime}{ }_{i}=b_{i}, \boldsymbol{i} \in \boldsymbol{S}_{\mathbf{1}} \cup \boldsymbol{S}_{\mathbf{2}}$
${b^{\prime}}_{i}=0, i \in S_{3} \quad, i=1,2, \cdots \cdots, m$,
$j=1,2, \cdots \cdots, n$
and form a new TMTP-MC.

## Step-3:

Calculate total minimum $\sup S_{\text {min }}$ andtotal minimum demand $D_{\text {min }}$
$S_{\text {min }}=\sum a_{i}, i=1,2, \cdots \cdots, m$
$D_{\min }=\sum b_{j}, j=1,2, \cdots \cdots, n$
Step-4: According to our assumption the supply and demand constraints will be
Supply Constraints

$$
\begin{gathered}
\sum_{j \in T} x_{i j}=a_{i}, \quad i \in S_{1} \\
\sum_{j \in T} x_{i j}=D_{\min }, \quad i \in S_{2} \\
\sum_{j \in T} x_{i j}=a_{i}, \quad i \in S_{3}
\end{gathered}
$$

and
Demand Constraints

$$
\begin{gathered}
\sum_{i \in S} x_{i j}=b_{j}, \quad j \in T_{1} \\
\sum_{i \in S} x_{i j}=S_{\min }, \quad j \in T_{2} \\
\sum_{i \in S} x_{i j}=b_{j}, \quad j \in T_{3}
\end{gathered}
$$

Finally, get the new supply $S_{\text {max }}^{\prime}$ and new demand $D^{\prime}{ }_{\text {max }}$.
$S_{\text {max }}^{\prime}=\sum a_{i}+\sum D_{\text {min }} \quad, i=1,2, \cdots \cdots, m$
$D^{\prime}{ }_{\text {max }}=\sum b_{j}+\sum S_{\text {min }}, j=1,2, \cdots \cdots, n$
Step-5: If necessarily make it balance adding zero row or column.
Step-6: Locate maximum time. Let us choose the row and column that pass the maxima time. From maximum time to minimum time allocate maximum possible allotment. If two or more minimum time exist choose which cell get maximum amount. Then choose next minimum time either same row or column and allocate as amount as possible and crossed out the row and column.
Step-7: If all supply and demand exhausted, then stop.
Step-8: Find total maximum transported unit $\sum x_{i j}$ and minimum required time.

## 4. Numerical Illustration:

## Example 1

For an emergency situation such as an Earthquake, government opened five emergency relief collection booths in five corners (airport) of Turkey. They collect medicine/food/necessary accessories from three reputed donating countries such as U, S and A are situated at different distances from Turkey. Production capacity of $U$ is exactly 80 units, A at most 140 units and that of S has at least 120 units. Moreover, collection booth-1 having a limit of demand exact 11 -unit, booth-2 having a limit of demand at least 13 unit, booth-3 having a limit of demand exact 60 units, booth- 4 having a limit of demand at most 80 units and collection booth- 5 having a limit of demand at least 80 units. Unit transportation time from source to collection booth are given below by this matrix.
From this type of TP with mixed constraints and we conclude the maximun transported unit with minimum transportation time:

Step-1: Formulation of Time Minimization Transportation Problem with Mixed Constraints (TMTP-MC).

|  |  | Collection Booth |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{4}$ | $\mathrm{D}_{5}$ | Supply |
| 2 | U | 8 | 6 | 8 | 6 | 10 | $=80$ |
|  | S | 9 | 8 | 9 | 6 | 7 | $\geq 120$ |
|  | A | 11 | 10 | 8 | 13 | 9 | $\leq 140$ |

Step-2, Step-3 and Step-4:


## Step-5:

New Equivalence StandardTime Minimization Transportation Problemwith Mixed Constraints will be according to our method:

|  |  | Collection Booth |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Supply |  |  |  |  |  |  |  |
|  |  | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{4}$ | $\mathrm{D}_{5}$ |  |
|  | U | 8 | 6 | 8 | 6 | 10 | 80 |
|  | S | 9 | 8 | 9 | 6 | 7 | 220 |
|  | S | A | 11 | 10 | 8 | 13 | 9 |
| Demand |  | 40 | 200 | 60 | 80 | 200 | $440 / 580$ |

It is unbalance form of TP. Make it balance we get

|  |  | Collection Booth |  |  |  |  | Supply |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{4}$ | $\mathrm{D}_{5}$ |  |
| $\begin{aligned} & \text { E. } \\ & \text { 右 } \\ & \text { O } \end{aligned}$ | U | 8 | 6 | 8 | 6 | 10 | 80 |
|  | S | 9 | 8 | 9 | 6 | 7 | 220 |
|  | A | 11 | 10 | 8 | 13 | 9 | 140 |
|  |  | 0 | 0 | 0 | 0 | 0 | 140 |
| Demand |  | 40 | 200 | 60 | 80 | 200 | 580 |

## Step-6, Step-7 and Step-8:

According to our proposed rule in Step-6, Step-7, Step-8, finally we get the table below:

|  | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{4}$ | D | Sup. |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{O}_{1}$ | 8 | 80 |  |  |  |  |
|  |  | 6 | 8 | 6 | 10 | =80 |
| $\mathrm{O}_{2}$ | 409 | 120 |  |  | 60 |  |
|  |  | 8 | 9 | 6 | 7 | $\geq 120$ |
| $\mathrm{O}_{3}$ | 11 | 10 | 8 | 13 | 80 |  |
|  |  |  |  |  | 9 | $\leq 140$ |
| Dem. | =40 | $\geq 40$ | =60 | $\leq 80$ | $\geq 80$ | 440 |

Therefore, the solution for the given problem is $x_{12}=80, x_{21}=40, x_{22}=120, x_{25}=60, x_{33}=$ 60, and $x_{35}=80$ and the corresponding time of the cells are $t_{12}=6, t_{21}=9 t_{22}=8, t_{25}=7$, $t_{33}=8$, and $t_{35}=9$.
Therefore, the total transportation time required

$$
\begin{aligned}
& \boldsymbol{T}_{\mathbf{0}}=\boldsymbol{\operatorname { m a x }}\left\{t_{12}, \quad t_{21}, t_{22}, \quad t_{25}, t_{33}, t_{35}\right\} \\
= & \max \{6,9,8,7,8,9\}=9
\end{aligned}
$$

So, the total transported unit is 440 and required maximum time 9 .

## 5. Conclusions

When transporting urgent materials, the amount transported matters more than the transportation cost. We developed an algorithm that will shipped maximum unit in minimum time satisfying all supply and demand requirements while over passing the transportation cost. The method proposed in this paper is very simple and easy understandable. No significant paper has been published on this context yet, on the contrary, our
solution technique is efficient and justified as the used data in the paper real. It provides solution in a simple manner, and examination could be helpful for troughs in pursuing vital choices to managing the emergency situation.

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